

Initially, the control unit determines whether the user's financial account is up-to-date. If not, the user unit, and thus the user, are notified and all further operations by the control unit are interrupted until the outstanding account is cleared up.

The decrypted and decompressed text of the IP is written into the control unit buffer. When the IP processing is complete, the user unit is notified that the IP is ready for loading via the SCSI bus. Also, if the billing register is not full, the user unit is also notified.

The pay-per-package accounting system in the information distribution system according to the present invention is illustrated in FIG. 15. At periodic intervals, or when the billing register is nearly full, the control unit transfers the billing register contents to the user unit for transfer to the central accounting office. Prior to transfer, the control unit generates a verification block key and determines the billing register verification block to provide security for the billing register contents.

Once valid information is transmitted, the central accounting office invoices the user for the information selected, retrieved and used. For example, the central office can print and mail an itemized invoice, or it can transmit the invoice directly to the user unit via the telephone connection for printing at the user unit.

In the case where the information packages (IP's) are transmitted by broadcast, the total length of an IP header block is preferably 64 bytes. The first 12 bytes are a synchronization string to allow a receiver to synchronize to the broadcast data stream. The next 2 bytes are the block count of the IP (its length in units of 32 byte blocks, not including the header block). The next 12 bytes are a message key. The following 22 bytes contain the article identification and classification information. The next 8 to 12 bytes are a checksum for the IP described by the header. The last 8 to 4 bytes are a

It should be emphasized that the IP header format may be varied, as desired, in accordance with the system requirements.

The preferred byte allocation is shown in the following Table:

Broadcast IP Header

XXXX	
XXXX	synchronization block
XXXX	
XX	IP block count
XXXX	message key
XXXX	
XX	
XXXX	
XXXX	IP identification and
XXXX	classification information
XXXX	
XXXX	
XXXX	
XXXX	IP validation block
XXXX	
XXXX	IP header validation block

The IP according to the present invention thus has a simple 3 part structure: The header block, the compressed dictionary, and the text block. If H represents the header block, D represent data blocks in the dictionary, and T represents data blocks in the text portion of the IP, the IP can be diagrammed in the following manner. The simplest structure for the IP is the header, followed by the dictionary, followed by the text block:

If block transposition is used to disrupt the message structure and increase the cryptographic resistance of the system, this simple structure is made apparently random, although the header will still appear as the lead block:

In this case it is necessary to use the keys to determine the transpose mapping so that the blocks can be fed through the decryption and decompression processors in the correct order.

The accounting (billing) register structure is also quite simple. This register contains a list of user access codes (if any), the associated IP identification numbers, as well as the IP prices, if desired (although the central billing office can look the prices up in its central databank.) The advantage of providing IP prices at the user site is for user convenience in selecting each IP and to maintain a local billing total that registers the current total billing amount. It is unlikely that publishers will want any more information be provided to users. 12